

# Fully Observable + Multiple Agents

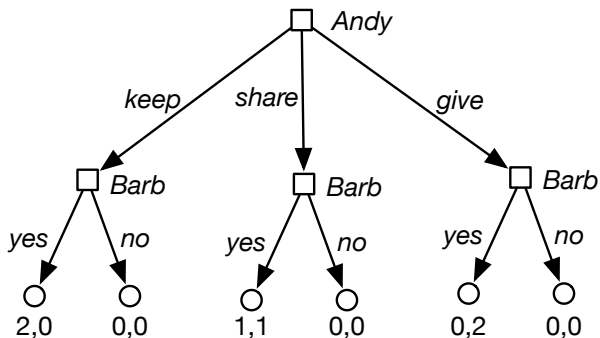
- If agents act sequentially and can observe the state before acting: **Perfect Information Games**.
- Can do dynamic programming or search:  
Each agent maximizes for itself.
- Multi-agent MDPs: value function for each agent.  
each agent maximizes its own value function.
- Multi-agent reinforcement learning: each agent has its own  $Q$  function.

# Fully-observable Game Tree Search

```
1: procedure GameTreeSearch(n)
2:   Inputs
3:     n a node in a game tree
4:   Output
5:     A pair: value for each agent for node n, path that gives
      this value
6:   if n is a leaf node then
7:     return {i : evaluate(i, n)}, None
8:   else if n is controlled by agent i then
9:     max :=  $-\infty$ 
10:    for each child c of n do
11:      score, path := GameTreeSearch(c)
12:      if score[i] > max then
13:        max := score[i]
14:        res := (score, c : path)
15:    return res
```

# Extensive Form of a Game

What happens with this game? Payoff is for *Andy*, *Barb*



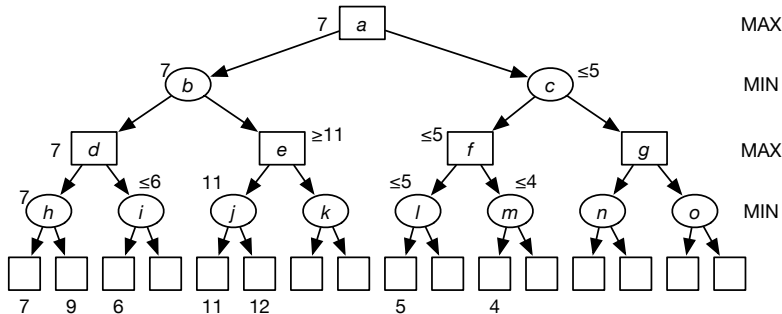
What if the 2,0 payoff was 1.9,0.1?

Should Barb be rational / predictable?

What should Andy do if Barb threatens to not do her best action?

# Pruning Dominated Strategies

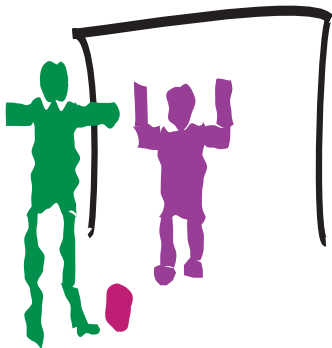
Special case: two person, competitive (zero sum) game. Utility for one agent is negative of utility of other agent.  $\implies$  minimax.



- square MAX nodes controlled by maximizing agent score
- round MIN nodes are controlled by a minimizing adversary
- leaves without a number do not need to be evaluated.

$\implies$   $\alpha$ - $\beta$  pruning.

# Partial Observability and Competition



		goalkeeper	
		left	right
kicker	left	0.6	0.2
	right	0.3	0.9

Probability of a goal.

- Each agent decides what to do without seeing the other agent's action.
- What should each agent do?

# Strategy Profiles

- Assume an  $n$ -player game in normal form
- A **strategy** for an agent is a probability distribution over the actions for this agent.
- A **strategy profile** is an assignment of a strategy to each agent.
- A strategy profile  $\sigma$  has a utility for each agent.  
Let  $utility(\sigma, i)$  be the utility of strategy profile  $\sigma$  for agent  $i$ .
- If  $\sigma$  is a strategy profile:  
 $\sigma_i$  is the strategy of agent  $i$  in  $\sigma$ ,  
 $\sigma_{-i}$  is the set of strategies of the other agents.  
Thus  $\sigma$  is  $\sigma_i\sigma_{-i}$

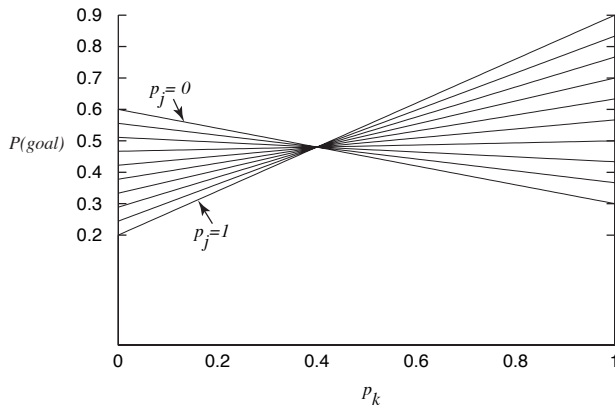
# Nash Equilibria

- $\sigma_i$  is a **best response** to  $\sigma_{-i}$  if for all other strategies  $\sigma'_i$  for agent  $i$ ,

$$utility(\sigma_i \sigma_{-i}, i) \geq utility(\sigma'_i \sigma_{-i}, i).$$

- A strategy profile  $\sigma$  is a **Nash equilibrium** if for each agent  $i$ , strategy  $\sigma_i$  is a best response to  $\sigma_{-i}$ . That is, a Nash equilibrium is a strategy profile such that no agent can do better by unilaterally deviating from that profile.
- Theorem [Nash, 1950] Every finite game has at least one Nash equilibrium.

# Stochastic Policies



		goalkeeper	
		left	right
kicker	left	0.6	0.2
	right	0.3	0.9

Probability of a goal.

$p_k$  is  $P(\text{kicker} = \text{right})$   
 $p_j$  is  $P(\text{goalkeeper} = \text{right})$



# Multiple Equilibria

Hawk-Dove Game:

		Agent 2	
		dove	hawk
Agent 1	dove	$R/2, R/2$	$0, R$
	hawk	$R, 0$	$-D, -D$

$D$  and  $R$  are both positive with  $D \gg R$ .

Just because you know the Nash equilibria doesn't mean you know what to do:

		Agent 2	
		shopping	football
Agent 1	shopping	2,1	0,0
	football	0,0	1,2

# Prisoner's Dilemma

Two strangers are in a game show. They each have the choice:

- Take \$100 for yourself
- Give \$1000 to the other player

This can be depicted as the payoff matrix:

		Player 2	
		take	give
Player 1	take	100,100	1100,0
	give	0,1100	1000,1000

# Tragedy of the Commons

Example:

- There are 100 agents.
- There is a common environment that is shared amongst all agents. Each agent has  $1/100$  of the shared environment.
- Each agent can choose to do an action that has a payoff of +10 but has a -100 payoff on the environment or do nothing with a zero payoff
- For each agent, doing the action has a payoff of  $10 - 100/100 = 9$
- If every agent does the action the total payoff is  $1000 - 10000 = -9000$

To compute a Nash equilibria for a game in strategic form:

- Eliminate dominated strategies
- Determine which actions will have non-zero probabilities. This is the **support set**.
- Determine the probability for the actions in the support set

# Eliminating Dominated Strategies

		Agent 2		
		$d_2$	$e_2$	$f_2$
Agent 1	$a_1$	3,5	5,1	1,2
	$b_1$	1,1	2,9	6,4
	$c_1$	2,6	4,7	0,8

- Can prune  $c_1$  because it is dominated by  $a_1$
- Can prune  $f_2$  because it is dominated by  $[0.5 : d_2, 0.5 : e_2]$
- Next prune  $b_1$  then  $e_2$
- Single Nash equilibrium is  $(a_1, d_2)$

Given a support set:

- Why would an agent will randomize between actions  $a_1 \dots a_k$ ?  
Actions  $a_1 \dots a_k$  have the same value for that agent given the strategies for the other agents.
- This forms a set of simultaneous equations where variables are probabilities of the actions
- A solution with all probabilities in range (0,1) is a Nash equilibrium.

Search over support sets to find a Nash equilibrium

## Example: computing Nash equilibrium

		goalkeeper	
		left	right
kicker	left	0.6	0.2
	right	0.3	0.9

Probability of a goal.

When would goalkeeper randomize?

Let  $p_k$  be probability the kicker will kick right.

$$\begin{aligned}P(\text{goal} \mid \text{jump left}) &= P(\text{goal} \mid \text{jump right}) \\p_k * 0.3 + (1 - p_k) * 0.6 &= p_k * 0.9 + (1 - p_k) * 0.2 \\0.6 - 0.2 &= (0.6 - 0.3 + 0.9 - 0.2) * p_k \\p_k &= 0.4\end{aligned}$$

Similarly for goal keeper:  $P(\text{jump right}) = 0.3$

Probability of a goal is:

$$(0.6 * 0.7) * 0.6 + (0.6 * 0.3) * 0.2 + (0.4 * 0.7) * 0.3 + (0.4 * 0.3) * 0.9 = 0.48$$



# Fictitious Play

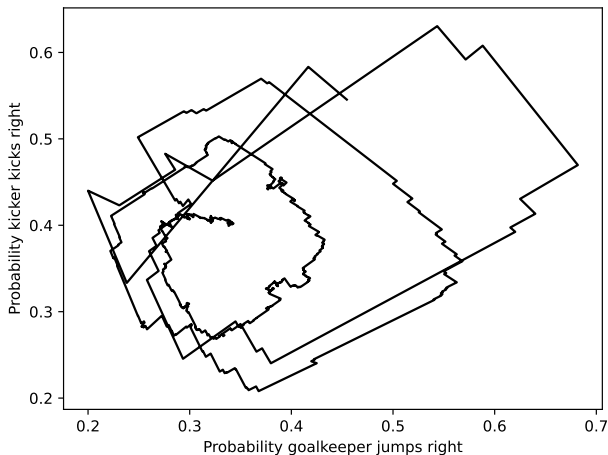
- Collect statistics of the other player.
- Assuming those statistics reflect the stochastic policy of the other agent, play a best response.
- Both players using fictitious play converges to a Nash equilibrium for many types of games (including two-player zero-sum games).
- If an opposing agent knew the exact strategy (whether learning or not) agent  $A$  was using, and could predict what agent  $A$  would do, it could exploit that knowledge.

```

1: controller Stochastic_policy_iteration( $S, A, \alpha, \gamma, q\_init, p\_init$ )
2:   Inputs
3:      $S$  is states,  $A$  is actions,  $\alpha$  is step size,  $\gamma$  discount
4:      $q\_init$  and  $p\_init$  ( $> 0$ ) are initial  $Q$  and  $P$  values
5:   Local
6:      $P[S, A]$  unnormalized  $P(a | s)$  ▷ Dirichlet
7:      $Q[S, A]$  estimate of value of doing  $A$  in state  $S$ 
8:      $P[s, a] := p\_init; Q[s, a] := q\_init$  for each  $s \in S$  and  $a \in A$ 
9:   observe state  $s$ ; select action  $a$  at random
10:  repeat
11:    do( $a$ )
12:    observe reward  $r$ , state  $s'$ 
13:    select action  $a'$  based on  $P[s', a'] / \sum_{a''} P[s', a'']$ 
14:     $Q[s, a] := Q[s, a] + \alpha * (r + \gamma * Q[s', a'] - Q[s, a])$ 
15:     $a\_best := \arg \max_a (Q[s, a])$ 
16:     $P[s, a\_best] = P[s, a\_best] + 1$ 
17:     $s := s'; a := a'$ 
18:  until termination

```

# Stochastic Policies



Repeated playing goal-kick game with single state ( $\alpha = 0.1$ ,  $\gamma = 0$ ,  $q\_init = 1$ ,  $p\_init = 5$ ).

ALPython: `masLearn.py`

**AlphaZero** – plays world-class chess, shogi, and Go. Improvement of program that beat Lee Sedol in 2016.

- implements **modified policy iteration**
- uses a deep neural network:
  - ▶ Input: the board position
  - ▶ Output: the value function and a stochastic policy
- To get a better estimate of the current state, it does stochastic simulation (**forward sampling**) of the rest of the game, using the stochastic policy with the upper confidence bound (UCB).
- This relies on a model to restart the search from any point.
- It was trained on self-play, playing itself for tens of millions of games.