At the end of the class you should be able to:

- justify why depth-bounded search is useful
- demonstrate how iterative-deepening works for a particular problem
- demonstrate how depth-first branch-and-bound works for a particular problem

- A bounded depth-first search takes a bound (cost or depth) and does not expand paths that exceed the bound.
  - explores part of the search graph
  - uses space linear in the depth of the search.
- How does this relate to other searches?
- How can this be extended to be complete?

## Which shaded goal will a depth-bounded search find first?



- Start with a bound b = 0.
- Do a bounded depth-first search with bound b
- If a solution is found return that solution
- Otherwise increment *b* and repeat.

- Start with a bound b = 0.
- Do a bounded depth-first search with bound b
- If a solution is found return that solution
- Otherwise increment *b* and repeat.
- This will find the same first solution as what other method?

- Start with a bound b = 0.
- Do a bounded depth-first search with bound b
- If a solution is found return that solution
- Otherwise increment *b* and repeat.
- This will find the same first solution as what other method?
- How much space is used?

- Start with a bound b = 0.
- Do a bounded depth-first search with bound b
- If a solution is found return that solution
- Otherwise increment *b* and repeat.
- This will find the same first solution as what other method?
- How much space is used?
- What happens if there is no path to a goal?

- Start with a bound b = 0.
- Do a bounded depth-first search with bound b
- If a solution is found return that solution
- Otherwise increment *b* and repeat.
- This will find the same first solution as what other method?
- How much space is used?
- What happens if there is no path to a goal?
- Surely recomputing paths is wasteful!!!

## Complexity with solution at depth k & branching factor b:

level	breadth-first	iterative deepening	# nodes
1	1	k	Ь
2	1	k-1	<i>b</i> <sup>2</sup>
k-1	1	2	$b^{k-1}$
k	1	1	b <sup>k</sup>
total			

## Complexity with solution at depth k & branching factor b:

level	breadth-first	iterative deepening	# nodes
1	1	k	Ь
2	1	k-1	<i>b</i> <sup>2</sup>
k-1	1	2	$b^{k-1}$
k	1	1	$b^k$
total	$\geq b^k$	$\leq b^k \left(rac{b}{b-1} ight)^2$	

< 🗆 )

- combines depth-first search with heuristic information.
- finds optimal solution.
- most useful when there are multiple solutions, and we want an optimal one.
- uses the space of depth-first search.

- Suppose *bound* is the cost of the lowest-cost path found to a goal so far.
- What if the search encounters a path p such that cost(p) + h(p) ≥ bound?

- Suppose *bound* is the cost of the lowest-cost path found to a goal so far.
- What if the search encounters a path p such that cost(p) + h(p) ≥ bound? p can be pruned.
- What can we do if a non-pruned path to a goal is found?

< 🗆 )

- Suppose *bound* is the cost of the lowest-cost path found to a goal so far.
- What if the search encounters a path p such that cost(p) + h(p) ≥ bound? p can be pruned.
- What can we do if a non-pruned path to a goal is found? bound can be set to the cost of p, and p can be remembered as the best solution so far.
- Why should this use a depth-first search?

- Suppose *bound* is the cost of the lowest-cost path found to a goal so far.
- What if the search encounters a path p such that cost(p) + h(p) ≥ bound? p can be pruned.
- What can we do if a non-pruned path to a goal is found? bound can be set to the cost of p, and p can be remembered as the best solution so far.
- Why should this use a depth-first search? Uses linear space.
- What can be guaranteed when the search completes?

- Suppose *bound* is the cost of the lowest-cost path found to a goal so far.
- What if the search encounters a path p such that cost(p) + h(p) ≥ bound? p can be pruned.
- What can we do if a non-pruned path to a goal is found? bound can be set to the cost of p, and p can be remembered as the best solution so far.
- Why should this use a depth-first search? Uses linear space.
- What can be guaranteed when the search completes? It has found an optimal solution.

- Suppose *bound* is the cost of the lowest-cost path found to a goal so far.
- What if the search encounters a path p such that cost(p) + h(p) ≥ bound? p can be pruned.
- What can we do if a non-pruned path to a goal is found? bound can be set to the cost of p, and p can be remembered as the best solution so far.
- Why should this use a depth-first search? Uses linear space.
- What can be guaranteed when the search completes? It has found an optimal solution.
- How should the bound be initialized?

- The bound can be initialized to  $\infty$ .
- The bound can be set to an estimate of the optimal path cost. After depth-first search terminates either:

< 🗆 )

- The bound can be initialized to  $\infty$ .
- The bound can be set to an estimate of the optimal path cost. After depth-first search terminates either:
  - A solution was found.
  - No solution was found, and no path was pruned
  - No solution was found, and a path was pruned.

## Which shaded goals will be best solutions so far?



(□)