

# Single agent or multiple agents

- Many domains are characterized by multiple agents rather than a single agent.
- **Game theory** studies what agents should do in a multi-agent setting.
- Agents can be cooperative, competitive or somewhere in between.
- Agents that are strategic can't be modeled as nature.

# Multi-agent framework

- Each agent can have its own utilities.
- Agents select actions autonomously.
- Agents can have different information.
- The outcome can depend on the actions of all of the agents.
- Each agent's value depends on the outcome.

# Fully Observable + Multiple Agents

- If agents act sequentially and can observe the state before acting: **Perfect Information Games**.
- Can do dynamic programming or search:  
Each agent maximizes for itself.
- Multi-agent MDPs: value function for each agent.  
each agent maximizes its own value function.
- Multi-agent reinforcement learning: each agent has its own  $Q$  function.
- Two person, competitive (zero sum)  $\implies$  minimax.

# Normal Form of a Game

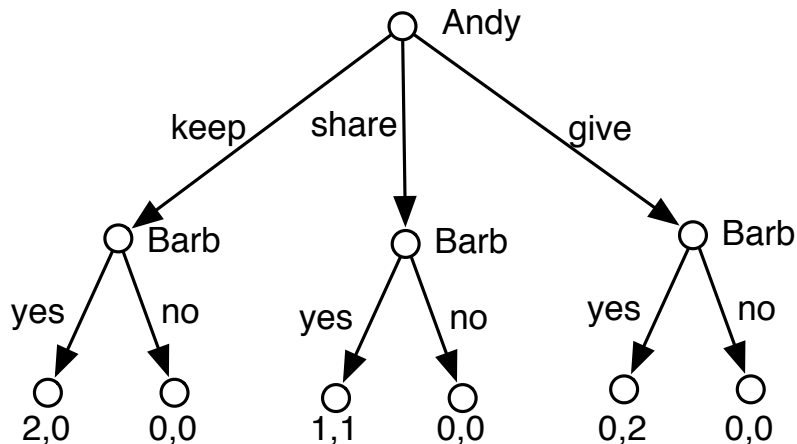
The **strategic form of a game** or **normal-form game**:

- a finite set  $I$  of agents,  $\{1, \dots, n\}$ .
- a set of actions  $A_i$  for each agent  $i \in I$ .  
An **action profile**  $\sigma$  is a tuple  $\langle a_1, \dots, a_n \rangle$ , means agent  $i$  carries out  $a_i$ .
- a utility function  $utility(\sigma, i)$  for action profile  $\sigma$  and agent  $i \in I$ , gives the expected utility for agent  $i$  when all agents follow action profile  $\sigma$ .

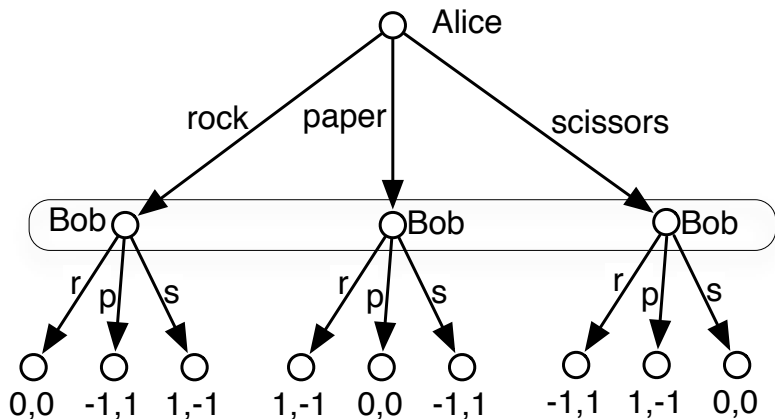
# Rock-Paper-Scissors

		Bob		
		<i>rock</i>	<i>paper</i>	<i>scissors</i>
Alice	<i>rock</i>	0, 0	-1, 1	1, -1
	<i>paper</i>	1, -1	0, 0	-1, 1
	<i>scissors</i>	-1, 1	1, -1	0, 0

# Extensive Form of a Game

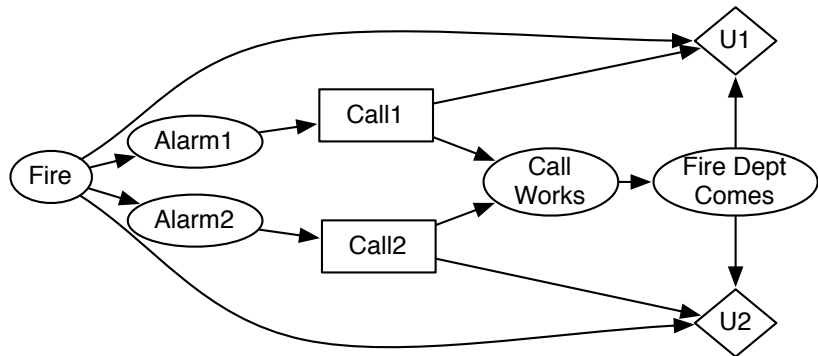


# Extensive Form of an imperfect-information Game



Bob cannot distinguish the nodes in an **information set**.

# Multiagent Decision Networks



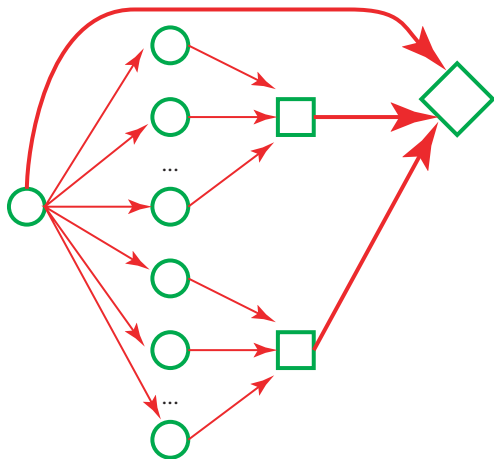
Value node for each agent.

Each decision node is owned by an agent.

Utility for each agent.



# Multiple Agents, shared value



# Complexity of Multi-agent decision theory

- It can be exponentially harder to find optimal multi-agent policy even with a shared values.
- **Why?** Because dynamic programming doesn't work:
  - ▶ If a decision node has  $n$  binary parents, dynamic programming lets us solve  $2^n$  decision problems.
  - ▶ This is much better than  $d^{2^n}$  policies (where  $d$  is the number of decision alternatives).
- Multiple agents with shared values is equivalent to having a single forgetful agent.

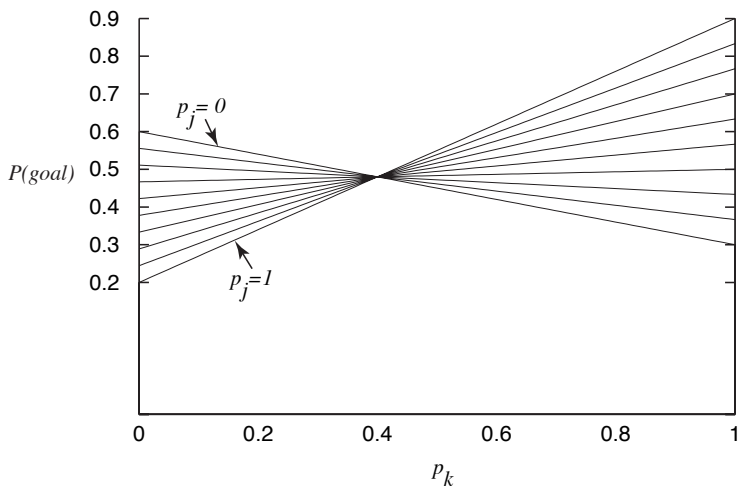
# Partial Observability and Competition



		goalie	
		left	right
kicker	left	0.6	0.2
	right	0.3	0.9

Probability of a goal.

# Stochastic Policies



# Strategy Profiles

- Assume a general  $n$ -player game,
- A **strategy** for an agent is a probability distribution over the actions for this agent.
- A **strategy profile** is an assignment of a strategy to each agent.
- A strategy profile  $\sigma$  has a utility for each agent. Let  $utility(\sigma, i)$  be the utility of strategy profile  $\sigma$  for agent  $i$ .
- If  $\sigma$  is a strategy profile:  
 $\sigma_i$  is the strategy of agent  $i$  in  $\sigma$ ,  
 $\sigma_{-i}$  is the set of strategies of the other agents.  
Thus  $\sigma$  is  $\sigma_i\sigma_{-i}$

# Nash Equilibria

- $\sigma_i$  is a **best response** to  $\sigma_{-i}$  if for all other strategies  $\sigma'_i$  for agent  $i$ ,

$$utility(\sigma_i \sigma_{-i}, i) \geq utility(\sigma'_i \sigma_{-i}, i).$$

- A strategy profile  $\sigma$  is a **Nash equilibrium** if for each agent  $i$ , strategy  $\sigma_i$  is a best response to  $\sigma_{-i}$ . That is, a Nash equilibrium is a strategy profile such that no agent can be better by unilaterally deviating from that profile.
- Theorem [Nash, 1950] Every finite game has at least one Nash equilibrium.

# Multiple Equilibria

Hawk-Dove Game:

		Agent 2	
		dove	hawk
Agent 1	dove	$R/2, R/2$	$0, R$
	hawk	$R, 0$	$-D, -D$

$D$  and  $R$  are both positive with  $D \gg R$ .

Just because you know the Nash equilibria doesn't mean you know what to do:

		Agent 2	
		shopping	football
Agent 1	shopping	2,1	0,0
	football	0,0	1,2



# Prisoner's Dilemma

Two strangers are in a game show. They each have the choice:

- Take \$100 for yourself
- Give \$1000 to the other player

This can be depicted as the payoff matrix:

		Player 2	
		take	give
Player 1	take	100,100	1100,0
	give	0,1100	1000,1000

# Tragedy of the Commons

Example:

- There are 100 agents.
- There is a common environment that is shared amongst all agents. Each agent has  $1/100$  of the shared environment.
- Each agent can choose to do an action that has a payoff of +10 but has a -100 payoff on the environment or do nothing with a zero payoff

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- If every agent does the action the total payoff is  $1000 - 10000 = -9000$

# Computing Nash Equilibria

To compute a Nash equilibria for a game in strategic form:

- Eliminate dominated strategies
- Determine which actions will have non-zero probabilities. This is the **support set**.
- Determine the probability for the actions in the support set

# Eliminating Dominated Strategies

		Agent 2		
		$d_2$	$e_2$	$f_2$
Agent 1	$a_1$	3,5	5,1	1,2
	$b_1$	1,1	2,9	6,4
	$c_1$	2,6	4,7	0,8

# Computing probabilities in randomized strategies

Given a support set:

- Why would an agent will randomize between actions  $a_1 \dots a_k$ ?



# Computing probabilities in randomized strategies

Given a support set:

- Why would an agent will randomize between actions  $a_1 \dots a_k$ ? Actions  $a_1 \dots a_k$  have the same value for that agent given the strategies for the other agents.
- This forms a set of simultaneous equations where variables are probabilities of the actions
- If there is a solution with all the probabilities in range  $(0,1)$  this is a Nash equilibrium.

Search over support sets to find a Nash equilibrium

# Learning to Coordinate

- Each agent maintains  $P[A]$  a probability distribution over actions.
- Each agent maintains  $Q[A]$  an estimate of value of doing  $A$  given policy of other agents.
- Repeat:
  - ▶ select action  $a$  using distribution  $P$ ,
  - ▶ do  $a$  and observe payoff
  - ▶ update  $Q$ :

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  - ▶ do  $a$  and observe payoff
  - ▶ update  $Q$ :  $Q[a] \leftarrow Q[a] + \alpha(\text{payoff} - Q[a])$
  - ▶ incremented probability of best action by  $\delta$ .
  - ▶ decremented probability of other actions